

Substructural Type Systems

CS 152 (Spring 2022)

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Today, we will learn about

- ▶ Substructural type systems
 - ▶ Natural deduction inference rules
 - ▶ Structural inference rules
 - ▶ Linear lambda calculus

- ▶ Rust

Natural deduction

- ▶ *Natural deduction* is a kind of proof calculus that can be used to formalize mathematical logic
 - ▶ Meant to be the natural way to reason about truth!
- ▶ $A_1, \dots, A_n \vdash B$ means whenever formulas A_1 to A_n are true, then formula B is true
- ▶ E.g., $p, \neg q \vdash q \Rightarrow (p \Rightarrow r)$
- ▶ Can define inference rules for the logic, e.g.,

$$\frac{\Gamma \vdash A \wedge B}{\Gamma \vdash A}$$

Structural inference rules

- ▶ *Structural inference rules* manipulate the assumptions (i.e., formulas to the left of \vdash)
- ▶ Allow us to treat list of formulas like a set.

Structural inference rule (1/3)

$$\text{EXCHANGE} \frac{\Gamma, A, B, \Delta \vdash C}{\Gamma, B, A, \Delta \vdash C}$$

Structural inference rule (2/3)

$$\text{CONTRACTION} \frac{\Gamma, A, A, \Delta \vdash B}{\Gamma, A, \Delta \vdash B}$$

Structural inference rule (3/3)

$$\text{WEAKENING} \frac{\Gamma, \Delta \vdash B}{\Gamma, A, \Delta \vdash B}$$

Natural deduction inference rules

- ▶ Additional inference rules needed, e.g., inference rules for propositional logic:

$$\frac{}{A \vdash A}$$

$$\frac{\Gamma, B \vdash A}{\Gamma \vdash B \Rightarrow A}$$

$$\frac{\Gamma \vdash B \Rightarrow A \quad \Delta \vdash B}{\Gamma, \Delta \vdash A}$$

$$\frac{\Gamma \vdash A \quad \Delta \vdash B}{\Gamma, \Delta \vdash A \wedge B}$$

$$\frac{\Gamma \vdash A \wedge B}{\Gamma \vdash A}$$

$$\frac{\Gamma \vdash A \wedge B}{\Gamma \vdash B}$$

Substructural logics

- ▶ If we drop any structural inference rule, we have a *substructural logic*

Substructural logics: linear logic

- ▶ Keep Exchange but drop Weakening and Contraction: **linear logic**
 - ▶ Every assumption must be used exactly once

$$\text{EXCHANGE} \frac{\Gamma, A, B, \Delta \vdash C}{\Gamma, B, A, \Delta \vdash C}$$

$$\text{CONTRACTION} \frac{\Gamma, A, A, \Delta \vdash B}{\Gamma, A, \Delta \vdash B}$$

$$\text{WEAKENING} \frac{\Gamma, \Delta \vdash B}{\Gamma, A, \Delta \vdash B}$$

Substructural logics: affine logic

- ▶ Keep Exchange and Weakening but
Contraction: **affine logic**
 - ▶ Every assumption must be used at most once

$$\text{EXCHANGE} \frac{\Gamma, A, B, \Delta \vdash C}{\Gamma, B, A, \Delta \vdash C}$$

$$\text{CONTRACTION} \frac{\Gamma, A, A, \Delta \vdash B}{\Gamma, A, \Delta \vdash B}$$

$$\text{WEAKENING} \frac{\Gamma, \Delta \vdash B}{\Gamma, A, \Delta \vdash B}$$

Curry-Howard Correspondence

Curry-Howard Correspondence

► Exchange

$$\frac{\Gamma, x:\tau_1, y:\tau_2, \Delta \vdash e:\tau}{\Gamma, y:\tau_2, x:\tau_1, \Delta \vdash e:\tau} \quad \frac{\Gamma, A, B, \Delta \vdash C}{\Gamma, B, A, \Delta \vdash C}$$

► Contraction

$$\frac{\Gamma, x:\tau, x:\tau, \Delta \vdash e:\tau'}{\Gamma, x:\tau, \Delta \vdash e:\tau'} \quad \frac{\Gamma, A, A, \Delta \vdash B}{\Gamma, A, \Delta \vdash B}$$

► Weakening

$$\frac{\Gamma, \Delta \vdash e:\tau}{\Gamma, x:\tau', \Delta \vdash e:\tau} \quad x \notin \Gamma, \Delta \quad \frac{\Gamma, \Delta \vdash B}{\Gamma, A, \Delta \vdash B}$$

Linear type system

- ▶ Every variable is used exactly once
- ▶ Linear type systems drop Contraction and Weakening (but keep Exchange)

$$\text{EXCHANGE} \frac{\Gamma, x:\tau_1, y:\tau_2, \Delta \vdash e:\tau}{\Gamma, y:\tau_2, x:\tau_1, \Delta \vdash e:\tau}$$

$$\text{CONTRACTION} \frac{\Gamma, x:\tau, x:\tau, \Delta \vdash e:\tau'}{\Gamma, x:\tau, \Delta \vdash e:\tau'}$$

$$\text{WEAKENING} \frac{\Gamma, \Delta \vdash e:\tau}{\Gamma, x:\tau', \Delta \vdash e:\tau} \quad x \text{ not in } \Gamma, \Delta$$

Affine type system

- ▶ Every variable is used at most once
- ▶ Affine type systems drop Contraction (but keep Exchange and Weakening)

$$\text{EXCHANGE} \frac{\Gamma, x:\tau_1, y:\tau_2, \Delta \vdash e:\tau}{\Gamma, y:\tau_2, x:\tau_1, \Delta \vdash e:\tau}$$

$$\text{CONTRACTION} \frac{\Gamma, x:\tau, x:\tau, \Delta \vdash e:\tau'}{\Gamma, x:\tau, \Delta \vdash e:\tau'}$$

$$\text{WEAKENING} \frac{\Gamma, \Delta \vdash e:\tau}{\Gamma, x:\tau', \Delta \vdash e:\tau} \quad x \text{ not in } \Gamma, \Delta$$

Relevant type system

- ▶ Every variable is used at least once
- ▶ Relevant type systems drop Weakening (but keep Contraction and Exchange)

$$\text{EXCHANGE} \frac{\Gamma, x:\tau_1, y:\tau_2, \Delta \vdash e:\tau}{\Gamma, y:\tau_2, x:\tau_1, \Delta \vdash e:\tau}$$

$$\text{CONTRACTION} \frac{\Gamma, x:\tau, x:\tau, \Delta \vdash e:\tau'}{\Gamma, x:\tau, \Delta \vdash e:\tau'}$$

$$\text{WEAKENING} \frac{\Gamma, \Delta \vdash e:\tau}{\Gamma, x:\tau', \Delta \vdash e:\tau} \quad x \text{ not in } \Gamma, \Delta$$

Ordered type system

- ▶ Every variable is used exactly once, in order
- ▶ Ordered type systems drop Weakening, Contraction, and Exchange

$$\text{EXCHANGE} \frac{\Gamma, x:\tau_1, y:\tau_2, \Delta \vdash e:\tau}{\Gamma, y:\tau_2, x:\tau_1, \Delta \vdash e:\tau}$$

$$\text{CONTRACTION} \frac{\Gamma, x:\tau, x:\tau, \Delta \vdash e:\tau'}{\Gamma, x:\tau, \Delta \vdash e:\tau'}$$

$$\text{WEAKENING} \frac{\Gamma, \Delta \vdash e:\tau}{\Gamma, x:\tau', \Delta \vdash e:\tau} \quad x \text{ not in } \Gamma, \Delta$$

Linear lambda calculus

- ▶ Explore linear type system in lambda calculus
- ▶ Type system will track use of objects
- ▶ A *linear object* must be used exactly once (and implementation could, e.g., deallocate object after use)
- ▶ Will also have unrestricted objects that can be used many times

Syntax

$$q ::= \text{lin} \mid \text{un}$$
$$e ::= x \mid q \ b \mid q \ (e_1, e_2) \mid q \ \lambda x:\tau. e \mid e_1 \ e_2 \\ \mid \text{if } e_1 \text{ then } e_2 \text{ else } e_3 \mid \text{split } e_1 \text{ as } x, y \text{ in } e_2$$
$$b \in \{\text{true}, \text{false}\}$$

Type system

$$\pi ::= \mathbf{bool} \mid \tau_1 \times \tau_2 \mid \tau_1 \rightarrow \tau_2$$
$$\tau ::= q \pi$$
$$\Gamma ::= \emptyset \mid \Gamma, x : \tau$$

Inference rules

- ▶ Maintain two invariants:
 1. linear variables are used exactly once on each control flow path
 2. unrestricted data structures may not contain linear data structures

Utility functions (1/2)

- ▶ Split context Γ into two pieces

$$\overline{\emptyset = \emptyset \circ \emptyset}$$

$$\Gamma = \Gamma_1 \circ \Gamma_2$$

$$\overline{\Gamma, x:\text{un } \pi = (\Gamma_1, x:\text{un } \pi) \circ (\Gamma_2, x:\text{un } \pi)}$$

$$\Gamma = \Gamma_1 \circ \Gamma_2$$

$$\overline{\Gamma, x:\text{lin } \pi = (\Gamma_1, x:\text{lin } \pi) \circ \Gamma_2}$$

$$\Gamma = \Gamma_1 \circ \Gamma_2$$

$$\overline{\Gamma, x:\text{lin } \pi = \Gamma_1 \circ (\Gamma_2, x:\text{lin } \pi)}$$

Utility functions (2/2)

- ▶ Determine whether type or context can be used in linear setting
 - ▶ $\text{un}(\tau)$ if and only if $\tau = \text{un } \pi$.
 - ▶ $\text{lin}(\tau)$ if and only if $\tau = \text{un } \pi$ or $\tau = \text{lin } \pi$.
 - ▶ $q(\Gamma)$ if and only if for all $(x:\tau) \in \Gamma$, we have $q(\tau)$.

Inference rules (1/2)

$$\frac{\text{un}(\Gamma_1, \Gamma_2)}{\Gamma_1, x:\tau, \Gamma_2 \vdash x:\tau}$$

$$\frac{\text{un}(\Gamma)}{\Gamma \vdash q \ b:q \ \mathbf{bool}}$$

$$\frac{\Gamma_1 \vdash e_1:q \ \mathbf{bool} \quad \Gamma_2 \vdash e_2:\tau \quad \Gamma_2 \vdash e_3:\tau}{\Gamma \vdash \text{if } e_1 \text{ then } e_2 \text{ else } e_3:\tau} \Gamma = \Gamma_1 \circ \Gamma_2$$

$$\frac{\Gamma_1 \vdash e_1:\tau_1 \quad \Gamma_2 \vdash e_2:\tau_2 \quad q(\tau_1) \quad q(\tau_2)}{\Gamma \vdash q(e_1, e_2):q(\tau_1, \tau_2)} \Gamma = \Gamma_1 \circ \Gamma_2$$

Inference rules (2/2)

$$\frac{\Gamma_1 \vdash e_1 : q (\tau_1 \times \tau_2) \quad \Gamma_2, x:\tau_1, y:\tau_2 \vdash e_2 : \tau}{\Gamma \vdash \text{split } e_1 \text{ as } x, y \text{ in } e_2 : \tau} \quad \Gamma = \Gamma_1 \circ \Gamma_2$$

$$\frac{q(\Gamma) \quad \Gamma, x:\tau \vdash e : \tau'}{\Gamma \vdash q \lambda x:\tau. e : q \tau \rightarrow \tau'}$$

$$\frac{\Gamma_1 \vdash e_1 : q \tau \rightarrow \tau' \quad \Gamma_2 \vdash e_2 : \tau}{\Gamma \vdash e_1 e_2 : \tau'} \quad \Gamma = \Gamma_1 \circ \Gamma_2$$

Example 1

lin λx : lin **bool**.

(lin λf : un (un **bool** \rightarrow lin **bool**). lin true)

(un λy : un **bool**. x)

Example 2

lin λx : lin **bool**.

(lin λf : un (un **bool** \rightarrow lin **bool**). lin (f (un true), f (un true)))
(un λy : un **bool**. x)

Operational semantics

- ▶ Use store-based semantics (to emphasize reclaiming memory)
- ▶ Type system ensures that a location is never accessed after it is freed.

$$p ::= b \mid \lambda x:\tau. e \mid (l_1, l_2)$$

$$v ::= q \ p$$

$$E ::= [\cdot] \mid \text{if } E \text{ then } e_2 \text{ else } e_3 \mid q (E, e) \mid q (l, E) \\ \mid \text{split } E \text{ as } x, y \text{ in } e \mid E \ e \mid l \ E$$

$$\langle e, \sigma \rangle \longrightarrow \langle e', \sigma' \rangle$$

$$\langle E[e], \sigma \rangle \longrightarrow \langle E[e'], \sigma' \rangle$$

$$\text{VAL} \frac{}{\langle v, \sigma \rangle \longrightarrow \langle l, \sigma[l \mapsto v] \rangle} \quad l \notin \text{dom}(\sigma)$$

$$\text{IF-TRUE} \frac{\sigma(l) = q \text{ true} \quad \sigma' = \begin{cases} \sigma & \text{if } q = \text{un} \\ \sigma \setminus l & \text{if } q = \text{lin} \end{cases}}{\langle \text{if } l \text{ then } e_1 \text{ else } e_2, \sigma \rangle \longrightarrow \langle e_1, \sigma' \rangle}$$

$$\text{IF-FALSE} \frac{\sigma(l) = q \text{ false} \quad \sigma' = \begin{cases} \sigma & \text{if } q = \text{un} \\ \sigma \setminus l & \text{if } q = \text{lin} \end{cases}}{\langle \text{if } l \text{ then } e_1 \text{ else } e_2, \sigma \rangle \longrightarrow \langle e_2, \sigma' \rangle}$$

$$\text{SPLIT} \frac{\sigma(l) = q (l_1, l_2) \quad \sigma' = \begin{cases} \sigma & \text{if } q = \text{un} \\ \sigma \setminus l & \text{if } q = \text{lin} \end{cases}}{\langle \text{split } l \text{ as } x, y \text{ in } e, \sigma \rangle \longrightarrow \langle e\{l_1/x\}\{l_2/y\}, \sigma' \rangle}$$

$$\text{APP} \frac{\sigma(l_1) = q \lambda x:\tau. e \quad \sigma' = \begin{cases} \sigma & \text{if } q = \text{un} \\ \sigma \setminus l_1 & \text{if } q = \text{lin} \end{cases}}{\langle l_1 \ l_2, \sigma \rangle \longrightarrow \langle e\{l_2/x\}, \sigma' \rangle}$$

Type Soundness

If $\vdash e:\tau$ and $\langle e, \emptyset \rangle \longrightarrow^* \langle e', \sigma \rangle$ then either

- ▶ e' is not stuck or
- ▶ $\exists l$ such that $e' = l$ and
 $\forall l' \in \text{dom}(\sigma)$ and $l' \notin \text{reachable}(l, \sigma)$, $\exists p$ such
that $\sigma(l') = \text{un } p$.