Harvard School of Engineering and Applied Sciences — CS 152: Programming Languages Induction; Small-step operational semantics; Large-step operational semantics Section and Practice Problems

Week 2: Tue Jan 31-Fri Feb 3, 2023

1 Induction

Let's inductively define a set of integers **Quux** with the following inference rules.

$$\text{RULE1} \frac{}{8 \in \mathbf{Quux}} \qquad \text{RULE2} \frac{}{5 \in \mathbf{Quux}} \qquad \text{RULE3} \frac{a \in \mathbf{Quux}}{c \in \mathbf{Quux}} c = a + b + 1$$

- (a) Of the rules above (i.e., RULE1, RULE2, and RULE3), which are axioms and which are inductive rules?
- (b) Give a derivation showing that 11 is in the set Quux.
- (c) Give a derivation showing that 20 is in the set **Quux**.
- (d) Write down the inductive reasoning principle for **Quux**. That is, if you wanted to prove that for some property P, for all $a \in \mathbf{Quux}$ we have P(a), what would you need to show? (See Lecture 3 §2.2 and §2.3.)
- (e) Prove that for all $a \in \mathbf{Quux}$, there exists $i \in \mathbb{Z}$ such that $a = 3 \times i 1$. Make sure that you follow the Recipe for Inductive Proofs! See Lecture 3 §2.5. What set are you inducting on? What is the property you are trying to prove? Go through each case.
- (f) Is 2 in the set **Quux**? If so, give a derivation proving it.

2 Small-step operational semantics

Consider the small-step operational semantics for the language of arithmetic expressions (Lecture 2). Let σ_0 be a store that maps all program variables to zero.

- (a) Show a derivation that $\langle 3 + (5 \times \mathsf{bar}), \sigma_0 \rangle \longrightarrow \langle 3 + (5 \times 0), \sigma_0 \rangle$.
- (b) What is the sequence of configurations that $\langle \text{foo} := 5; (\text{foo} + 2) \times 7, \sigma_0 \rangle$ steps to? (You don't need to show the derivations for each step, just show what configuration $\langle \text{foo} := 5; (\text{foo} + 2) \times 7, \sigma_0 \rangle$ steps to in one step, then two steps, then three steps, and so on, until you reach a final configuration.)
- (c) Find an integer n and store σ' such that $\langle ((6+(\mathsf{foo} := (\mathsf{bar} := 3; 5); 1+\mathsf{bar})) + \mathsf{bar}) \times \mathsf{foo}, \sigma_0 \rangle \longrightarrow^* \langle n, \sigma' \rangle$.
- (d) Is the relation → reflexive? Is it symmetric? Is it anti-symmetric? Is it transitive?
 (For each of these questions, if the answer is "no", what is a suitable counterexample? If any of the answers are "yes", think about how you would prove it.)

3 Large-step operational semantics

Consider the large-step operational semantics for the language of arithmetic expressions (Lecture 4). Let σ_0 be a store that maps all program variables to zero.

(a) Show a derivation that $\langle 3 + (5 \times \mathsf{bar}), \sigma_0 \rangle \downarrow \langle 3, \sigma_0 \rangle$.

- (b) Find an integer n and store σ' such that $\langle \mathsf{foo} := 5; (\mathsf{foo} + 2) \times 7, \sigma_0 \rangle \Downarrow \langle n, \sigma' \rangle$. If you have time and a big piece of paper, give the derivation of $\langle \mathsf{foo} := 5; (\mathsf{foo} + 2) \times 7, \sigma_0 \rangle \Downarrow \langle n, \sigma' \rangle$.
- (c) Is the relation \$\psi\$ reflexive? Is it symmetric? Is it anti-symmetric? Is it transitive?

 (For each of these questions, if the answer is "no", what is a suitable counterexample? If any of the answers are "yes", think about how you would prove it.)