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Abstract Interpretation: Fixpoints, widening, and narrowing

CS252r Spring 2011

*Slides from
Principles of Program Analysis
by Nielson, Nielson, and Hankin*

<http://www2.imm.dtu.dk/~riis/PPA/ppasup2004.html>

The need for fix-points

- Let L be complete lattice
- Suppose $f:L \rightarrow L$ is program analysis for some program construct p
 - i.e. $p \vdash l_1 \triangleright l_2$ where $f(l_1)=l_2$
 - monotonic function
- If p is recursive or iterative program construct, want to find **least fixed point** (lfp) of f .
 - Most precise lattice element representing analysis of executing p unbounded number of times

Fixed points

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Let $f : L \rightarrow L$ be a *monotone function* on a complete lattice $L = (L, \sqsubseteq, \sqcup, \sqcap, \perp, \top)$.

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l is a *fixed point* iff $f(l) = l$ $Fix(f) = \{l \mid f(l) = l\}$

f is *reductive* at l iff $f(l) \sqsubseteq l$ $Red(f) = \{l \mid f(l) \sqsubseteq l\}$

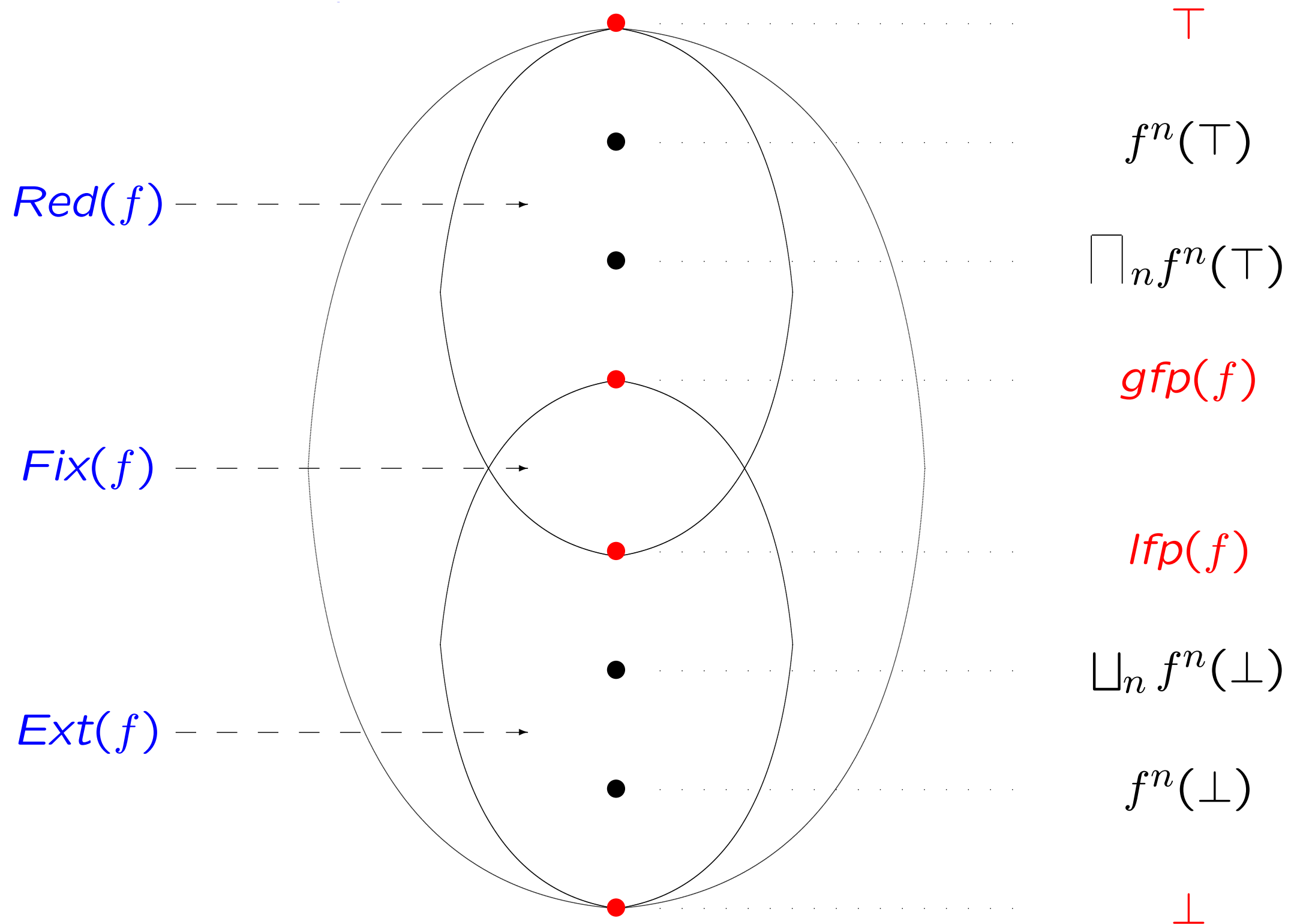
f is *extensive* at l iff $f(l) \sqsupseteq l$ $Ext(f) = \{l \mid f(l) \sqsupseteq l\}$

Tarski's Theorem ensures that

$$lfp(f) = \sqcap Fix(f) = \sqcap Red(f) \in Fix(f) \subseteq Red(f)$$

$$gfp(f) = \sqcup Fix(f) = \sqcup Ext(f) \in Fix(f) \subseteq Ext(f)$$

Fixed points of f



Need for approximation

- How do we find $\text{lfp}(f)$?
- Ideally use iterative sequence
 - $(f^n(\perp))_n = \perp, f(\perp), f(f(\perp)), \dots$
- But:
 - may not stabilize
 - if L doesn't meet ascending chain condition
 - least upper bound of $(f^n(\perp))_n$ may not equal $\text{lfp}(f)$
 - Why?
 - No guarantee f is continuous, and so Kleene's fixed-point theorem doesn't apply
- Need to approximate...

One possibility

- Start with \top and repeatedly apply f
 - i.e., $(f^n(\top))_n = \top, f(\top), f(f(\top)), \dots$
- Even if it doesn't stabilize, will always be a sound approximation
 - for all i we have $\text{lfp}(f) \sqsubseteq f^i(\top)$
 - Means that can stop when we run out of patience, and have sound approximation
- But in practice, too imprecise.

Widening operators

- Key idea: replace $(f^n(\perp))_n$ with sequence $(f_{\nabla}^n)_n$ such that
 - $(f_{\nabla}^n)_n$ guaranteed to stabilize with safe (upper) approximation of $\text{lfp}(f)$
- ∇ is a **widening operator**
 - An upper bound operator satisfying a finiteness condition

Upper bound operators

$\checkmark : L \times L \rightarrow L$ is an *upper bound operator* iff

$$l_1 \sqsubseteq l_1 \checkmark l_2 \sqsupseteq l_2$$

for all $l_1, l_2 \in L$.

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Let $(l_n)_n$ be a sequence of elements of L . Define the sequence $(l_n^{\checkmark})_n$ by:

$$l_n^{\checkmark} = \begin{cases} l_n & \text{if } n = 0 \\ l_{n-1}^{\checkmark} \checkmark l_n & \text{if } n > 0 \end{cases}$$

Fact: If $(l_n)_n$ is a sequence and \checkmark is an upper bound operator then $(l_n^{\checkmark})_n$ is an ascending chain; furthermore $l_n^{\checkmark} \sqsupseteq \sqcup\{l_0, l_1, \dots, l_n\}$ for all n .

Example

Let *int* be an arbitrary but fixed element of **Interval**.

An upper bound operator:

$$int_1 \sqcup^{int} int_2 = \begin{cases} int_1 \sqcup int_2 & \text{if } int_1 \sqsubseteq int \vee int_2 \sqsubseteq int_1 \\ [-\infty, \infty] & \text{otherwise} \end{cases}$$

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Example: $[1, 2] \sqcup^{[0,2]} [2, 3] = [1, 3]$ and $[2, 3] \sqcup^{[0,2]} [1, 2] = [-\infty, \infty]$.

Transformation of: $[0, 0], [1, 1], [2, 2], [3, 3], [4, 4], [5, 5], \dots$

If $int = [0, \infty]$: $[0, 0], [0, 1], [0, 2], [0, 3], [0, 4], [0, 5], \dots$

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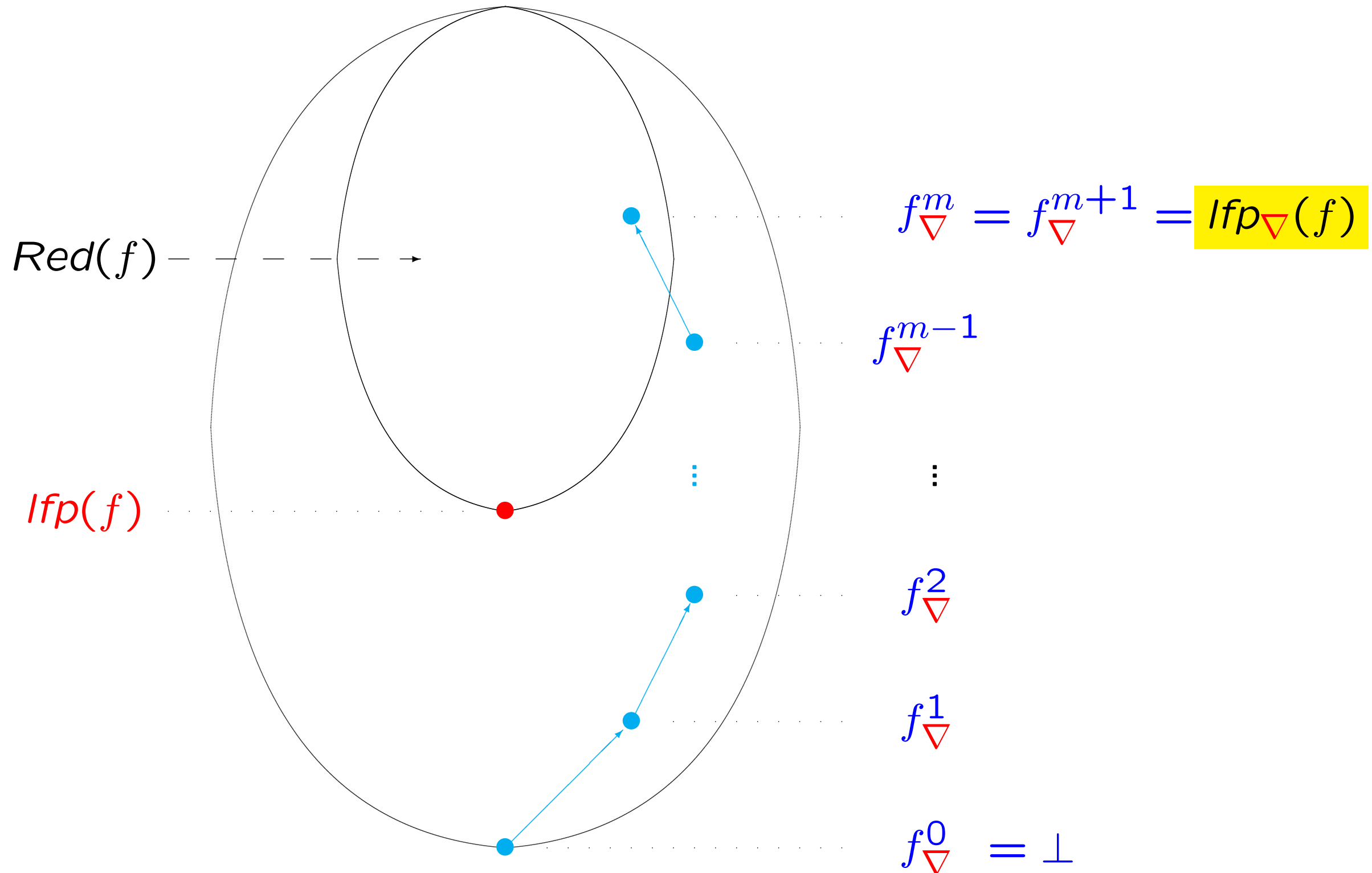
Widening operators

- Operator $\nabla : L \times L \rightarrow L$ is a **widening operator** iff
 - ∇ is an upper bound operator
 - for all ascending chains $(I_n)_n$ the ascending chain $(I_n^\nabla)_n$ eventually stabilizes
 - $I_n^\nabla = I_n$ if $n = 0$
 - $I_n^\nabla = I_{n-1}^\nabla \nabla I_n$ otherwise

Widening operators

- For monotonic function $f: L \rightarrow L$ and widening operator ∇ define $(f_{\nabla}^n)_n$ by
 - $f_{\nabla}^n = \perp$ if $n = 0$
 - $f_{\nabla}^n = f_{\nabla}^{n-1}$ if $n > 0$ and $f(f_{\nabla}^{n-1}) \sqsubseteq f_{\nabla}^{n-1}$
 - $f_{\nabla}^n = f_{\nabla}^{n-1} \nabla f(f_{\nabla}^{n-1})$ otherwise
- This is an ascending chain that eventually stabilizes
 - when $f(f_{\nabla}^m) \sqsubseteq f_{\nabla}^m$ for some m
- Tarski's Thm then gives $f_{\nabla}^m \sqsupseteq \text{lfp}(f)$

Diagrammatically



Example

Let K be a *finite* set of integers, e.g. the set of integers explicitly mentioned in a given program.

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We shall define a widening operator ∇ based on K .

Idea: $[z_1, z_2] \nabla [z_3, z_4]$ is

$$[\text{LB}(z_1, z_3) , \text{UB}(z_2, z_4)]$$

where

- $\text{LB}(z_1, z_3) \in \{z_1\} \cup K \cup \{-\infty\}$ is the best possible lower bound, and
- $\text{UB}(z_2, z_4) \in \{z_2\} \cup K \cup \{\infty\}$ is the best possible upper bound.

The effect: a change in any of the bounds of the interval $[z_1, z_2]$ can only take place finitely many times – corresponding to the cardinality of K .

Example

Let $z_i \in \mathbf{Z}' = \mathbf{Z} \cup \{-\infty, \infty\}$ and write:

$$\text{LB}_K(z_1, z_3) = \begin{cases} z_1 & \text{if } z_1 \leq z_3 \\ k & \text{if } z_3 < z_1 \wedge k = \max\{k \in K \mid k \leq z_3\} \\ -\infty & \text{if } z_3 < z_1 \wedge \forall k \in K : z_3 < k \end{cases}$$

$$\text{UB}_K(z_2, z_4) = \begin{cases} z_2 & \text{if } z_4 \leq z_2 \\ k & \text{if } z_2 < z_4 \wedge k = \min\{k \in K \mid z_4 \leq k\} \\ \infty & \text{if } z_2 < z_4 \wedge \forall k \in K : k < z_4 \end{cases}$$

$$int_1 \nabla int_2 = \begin{cases} \perp & \text{if } int_1 = int_2 = \perp \\ [\text{LB}_K(\text{inf}(int_1), \text{inf}(int_2)) , \text{UB}_K(\text{sup}(int_1), \text{sup}(int_2))] & \text{otherwise} \end{cases}$$

Example

Consider the ascending chain $(int_n)_n$

$$[0, 1], [0, 2], [0, 3], [0, 4], [0, 5], [0, 6], [0, 7], \dots$$

and assume that $K = \{3, 5\}$.

Then $(int_n^\nabla)_n$ is the chain

$$[0, 1], [0, 3], [0, 3], [0, 5], [0, 5], [0, \infty], [0, \infty], \dots$$

which eventually stabilises.

Defining widening operators

- Suppose we have two complete lattices, L and M , and a Galois connection (L, α, γ, M) between them
- One possibility: replace analysis $f:L \rightarrow L$ with analysis $g:M \rightarrow M$
 - Can **induce** g from f
 - But may reduce precision of analysis
- Another possibility
 - Use M just to ensure convergence of fixedpoints
 - Assume upper bound operator ∇_M for M
 - Define $l_1 \nabla_L l_2 = \gamma(\alpha(l_1) \nabla_M \alpha(l_2))$
 - ∇_L is widening operator if either
 - (i) M has no infinite ascending chains or
 - (ii) (L, α, γ, M) is Galois insertion and ∇_M is widening operator

Improving on $\text{lfp}_{\nabla}(f)$

- Widening gives upper approximation $\text{lfp}_{\nabla}(f)$ of $\text{lfp}(f)$
- But $f(\text{lfp}_{\nabla}(f)) \sqsubseteq \text{lfp}_{\nabla}(f)$ so we can improve approximation by considering sequence $(f^n(\text{lfp}_{\nabla}(f)))_n$
- For all i we have $\text{lfp}(f) \sqsubseteq f^i(\text{lfp}_{\nabla}(f)) \sqsubseteq \text{lfp}_{\nabla}(f)$
 - So can stop anytime with an upper approximation
- Defining a **narrowing operator** gives a way to describe when to stop

Narrowing operator

An operator $\Delta : L \times L \rightarrow L$ is a *narrowing operator* iff

- $l_2 \sqsubseteq l_1 \Rightarrow l_2 \sqsubseteq (l_1 \Delta l_2) \sqsubseteq l_1$ for all $l_1, l_2 \in L$, and
- for all descending chains $(l_n)_n$ the sequence $(l_n^\Delta)_n$ eventually stabilises.

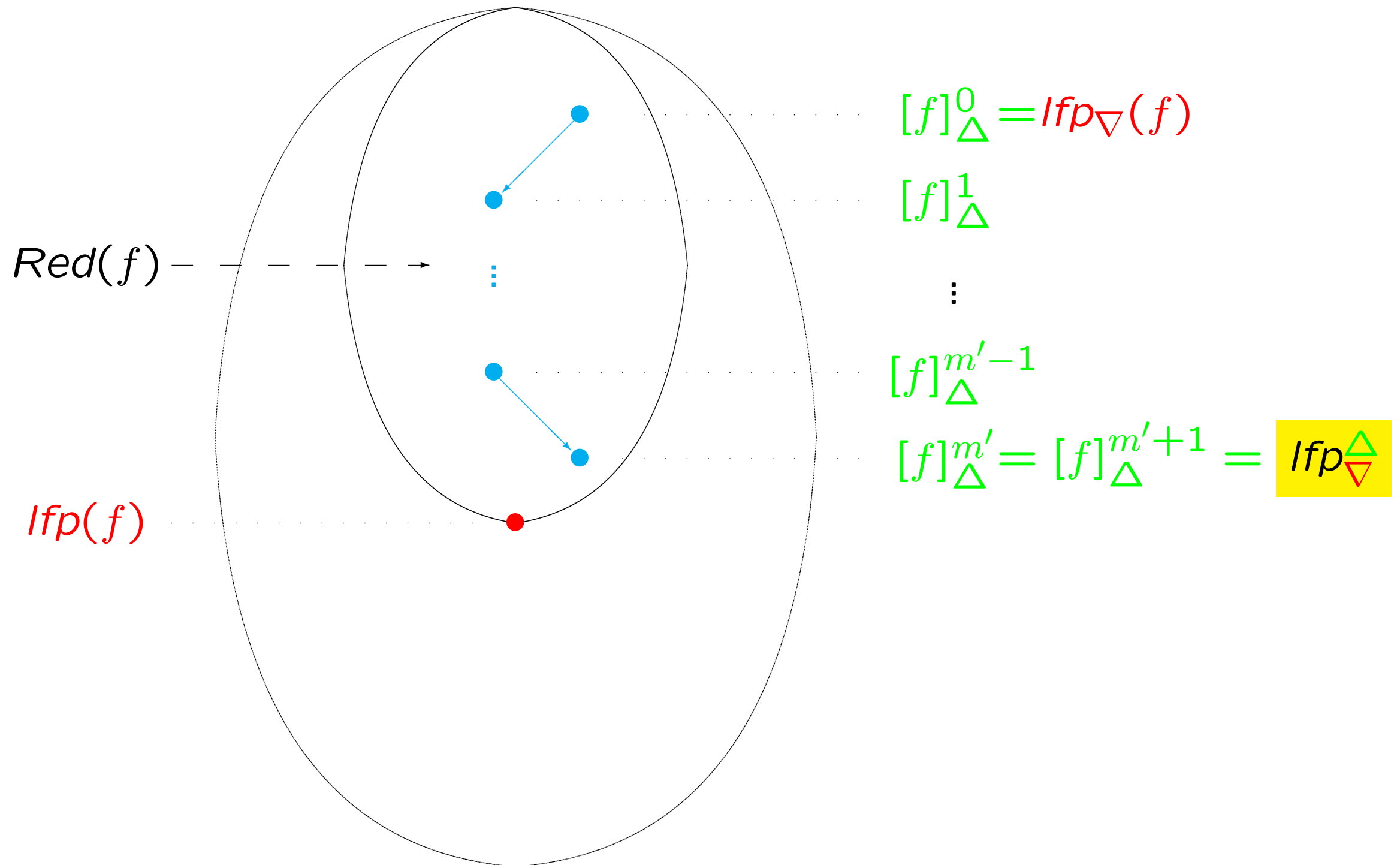
We construct the sequence $([f]_\Delta^n)_n$

$$[f]_\Delta^n = \begin{cases} \text{Ifp}_\nabla(f) & \text{if } n = 0 \\ [f]_\Delta^{n-1} \Delta f([f]_\Delta^{n-1}) & \text{if } n > 0 \end{cases}$$

One can show that:

- $([f]_\Delta^n)_n$ is a descending chain where all elements satisfy $\text{Ifp}(f) \sqsubseteq [f]_\Delta^n$
- the chain eventually stabilises so $[f]_\Delta^{m'} = [f]_\Delta^{m'+1}$ for some value m'

Diagrammatically



Example

The complete lattice (**Interval**, \sqsubseteq) has two kinds of infinite descending chains:

- those with elements of the form $[-\infty, z]$, $z \in \mathbf{Z}$
- those with elements of the form $[z, \infty]$, $z \in \mathbf{Z}$

Idea: Given some fixed non-negative number N the narrowing operator Δ_N will force an infinite descending chain

$$[z_1, \infty], [z_2, \infty], [z_3, \infty], \dots$$

(where $z_1 < z_2 < z_3 < \dots$) to stabilise when $z_i > N$

Similarly, for a descending chain with elements of the form $[-\infty, z_i]$ the narrowing operator will force it to stabilise when $z_i < -N$

Example

Define $\Delta = \Delta_N$ by

$$int_1 \Delta int_2 = \begin{cases} \perp & \text{if } int_1 = \perp \vee int_2 = \perp \\ [z_1, z_2] & \text{otherwise} \end{cases}$$

where

$$z_1 = \begin{cases} \inf(int_1) & \text{if } N < \inf(int_2) \wedge \sup(int_2) = \infty \\ \inf(int_2) & \text{otherwise} \end{cases}$$

$$z_2 = \begin{cases} \sup(int_1) & \text{if } \inf(int_2) = -\infty \wedge \sup(int_2) < -N \\ \sup(int_2) & \text{otherwise} \end{cases}$$

Consider the infinite descending chain $([n, \infty])_n$

$$[0, \infty], [1, \infty], [2, \infty], [3, \infty], [4, \infty], [5, \infty], \dots$$

and assume that $N = 3$.

Then the narrowing operator Δ_N will give the sequence $([n, \infty]^\Delta)_n$

$$[0, \infty], [1, \infty], [2, \infty], [3, \infty], [3, \infty], [3, \infty], \dots$$

Summary

- Given monotonic $f:L \rightarrow L$ where L is a lattice
- Approximating least fixed point of f accurately and quickly a key challenge of program analysis
- Widening operators
- Widening following by narrowing